

# A Coinductive Version of Milner's Proof System for Regular Expressions Modulo Bisimilarity

Clemens Grabmayer

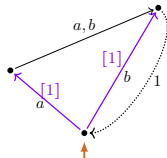
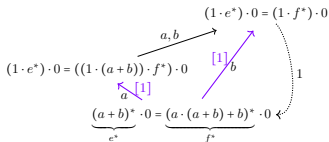
**G S** GRAN SASSO  
SCIENCE INSTITUTE

**S I** SCHOOL OF ADVANCED STUDIES  
Scuola Universitaria Superiore

Department of Computer Science  
L'Aquila, Italy

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# Process semantics $\llbracket \cdot \rrbracket_P$ of regular (star) expr's *(Milner, 1984)*

$0 \xrightarrow{P}$  deadlock  $\delta$ , no termination

$1 \xrightarrow{P}$  empty process  $\epsilon$ , then terminate

$a \xrightarrow{P}$  atomic action  $a$ , then terminate

$e + f \xrightarrow{P}$  alternative composition of  $\llbracket e \rrbracket_P$  and  $\llbracket f \rrbracket_P$

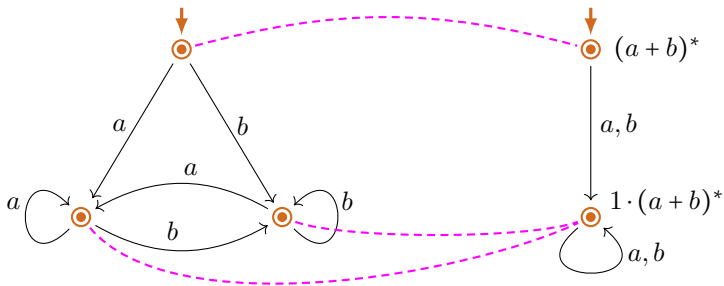
$e \cdot f \xrightarrow{P}$  sequential composition of  $\llbracket e \rrbracket_P$  and  $\llbracket f \rrbracket_P$

$e^* \xrightarrow{P}$  unbounded iteration of  $\llbracket e \rrbracket_P$ , option to terminate

$\llbracket e \rrbracket_P := [\mathbf{P}(e)]_{\Leftrightarrow}$  (bisimilarity equivalence class of process  $\mathbf{P}(e)$ )

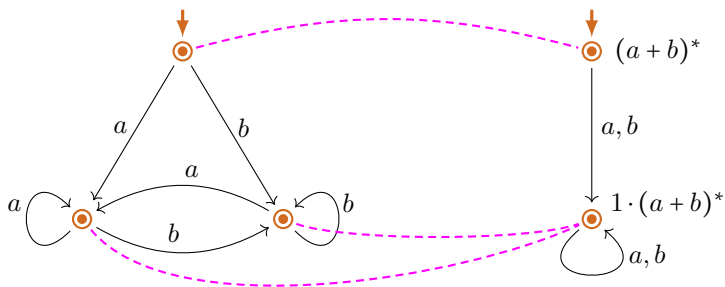
$:= [\mathcal{C}(e)]_{\Leftrightarrow}$  (bisimilarity equivalence class of chart  $\mathcal{C}(e)$ )

# Chart interpretation (example) (via TSS or *Antimirov's partial deriv's*)



$$\begin{aligned} \mathcal{C}((a^* \cdot b^*)^*) & \quad \Longleftrightarrow \quad \mathcal{C}((a + b)^*) \\ \llbracket (a^* \cdot b^*)^* \rrbracket_{\mathcal{P}} & \quad = \quad \llbracket (a + b)^* \rrbracket_{\mathcal{P}} \end{aligned}$$

# Chart interpretation (example) (via TSS or *Antimirov's partial deriv's*)



$$\begin{array}{ccc}
 \mathcal{C}((a^* \cdot b^*)^*) & \stackrel{\leftrightarrow}{=} & \mathcal{C}((a+b)^*) \\
 (a^* \cdot b^*)^* & \stackrel{=}{=} \llbracket \cdot \rrbracket_P & (a+b)^*
 \end{array}$$

# Milner's proof system $Mil = Mil^- + RSP^*$

## Axioms:

$$(assoc(+)) \quad (e + f) + g = e + (f + g)$$

$$(neutr(+)) \quad e + 0 = e$$

$$(comm(+)) \quad e + f = f + e$$

$$(idempot(+)) \quad e + e = e$$

$$(assoc(\cdot)) \quad (e \cdot f) \cdot g = e \cdot (f \cdot g)$$

$$(r-distr(+, \cdot)) \quad (e + f) \cdot g = e \cdot g + f \cdot g$$

$$(id_l(\cdot)) \quad 1 \cdot e = e$$

$$(id_r(\cdot)) \quad e \cdot 1 = e$$

$$(deadlock) \quad 0 \cdot e = 0$$

$$(rec(*)) \quad e^* = 1 + e \cdot e^*$$

$$(trm-body(*)) \quad e^* = (1 + e)^*$$

## Inference rules: equational logic plus

$$\frac{e = f \cdot e + g}{e = f^* \cdot g} RSP^* \quad (\text{if } f \downarrow)$$

# Milner's axiomatization question

## Question (*Milner, 1984*)

Is Milner's system  $\text{Mil} = \text{Mil}^- + \text{RSP}^*$  **complete** for bisimilarity of process interpretations of regular expressions?

$$\forall e, f \text{ reg. expr's } \left( \vdash_{\text{Mil}} e = f \begin{array}{c} \xleftarrow{\text{complete?}} \\ \xrightarrow{\text{sound}} \end{array} e = \llbracket \cdot \rrbracket_P f \right) ?$$

"Yes" for restrictions to subclasses: *Zantema/Fokkink (1994)*, *Fokkink (1996)*, *Corradini, De Nicola, Labella (2002)*, *G/Fokkink (2020)*.

## Proposition (*G, CMCS 2006*)

The system  $\text{Mil}^- + \text{USP}$ , where:

**USP**: **unique solvability** of guarded, linear systems of equations, is (sound and) **complete**.

## Question (investigated here)

How can the derivational power be characterized that the **fixed-point rule RSP\*** adds to the purely equational part  $\text{Mil}^-$  of Milner's system?

# Answer developed

We use:

- ▶ the loop existence and elimination property (**LEE**) of charts
  - ▶ implies expressibility by a star expression
  - ▶ led to completeness result for **1-free star expressions** (G/Fokkink, 2020)

We introduce:

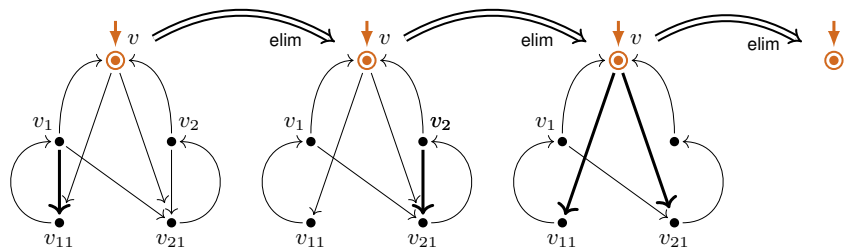
- ▶ a coinductive version **cMil** = (**Mil**<sup>-</sup> + **LCoProof**) of **Mil** = (**Mil**<sup>-</sup> + **RSP\***) based on **LEE-witnessed** coinductive proofs over **Mil**<sup>-</sup>.

We construct / obtain:

- ▶ a proof transformation: **Mil**  $\mapsto$  **cMil**, (**RSP\*** inst's  $\mapsto$  **LCoProof** inst's),
- ▶ a proof transformation: **Mil**  $\longleftarrow$  **cMil**, (bottom-up extraction procedure),
- ▶ theorem equivalence **Mil**  $\sim$  **cMil** :

$$\vdash_{\text{Mil}} e = f \iff \vdash_{\text{cMil}} e = f.$$

## LEE



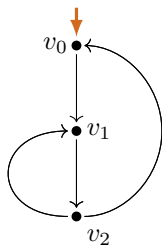


# Loop charts (interpretations of innermost iterations in 1-free expressions)

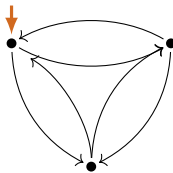
## Definition

A chart is a **loop chart** if:

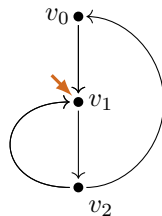
- (L1) There is an infinite path from the **start vertex**.
- (L2) Every infinite path from the **start vertex** returns to it.
- (L3) Immediate termination is **only** possible at the **start vertex**.



(L1), ~~(L2)~~



(L1), (L2), ~~(L3)~~

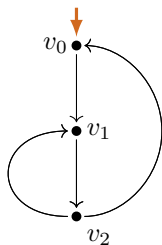


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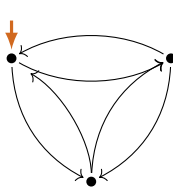
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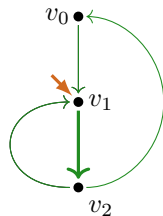
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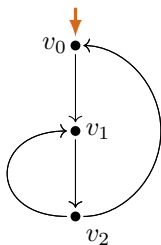


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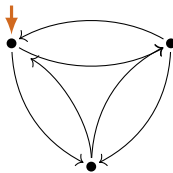
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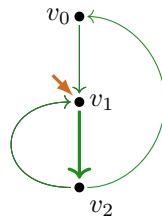
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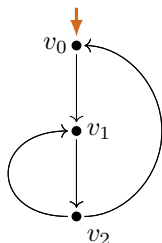
loop chart

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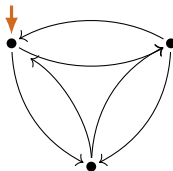
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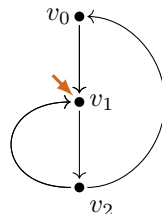
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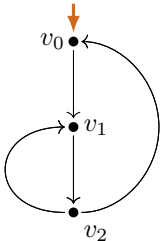
loop chart

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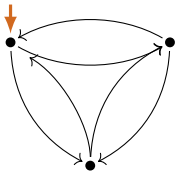
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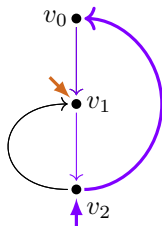
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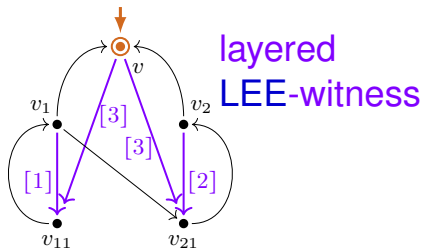
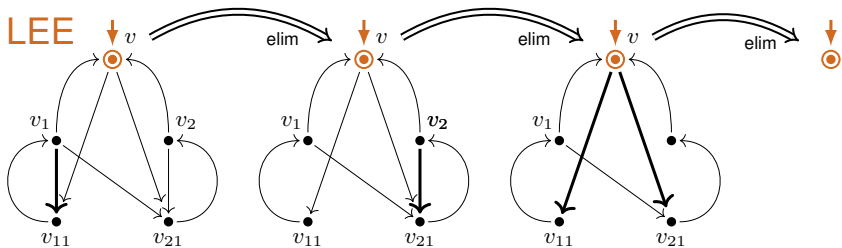


(L1), (L2), ~~(L3)~~



loop subchart

# LEE, and LLEE-witness

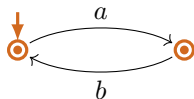


# LEE

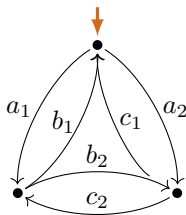
## Definition

A chart  $\mathcal{C}$  satisfies **LEE** (*loop existence and elimination*) if:

$$\exists \mathcal{C}_0 \left( \mathcal{C} \xRightarrow{*}_{\text{elim}} \mathcal{C}_0 \not\Rightarrow_{\text{elim}} \mathcal{C}_0 \wedge \mathcal{C}_0 \text{ permits no infinite path} \right).$$

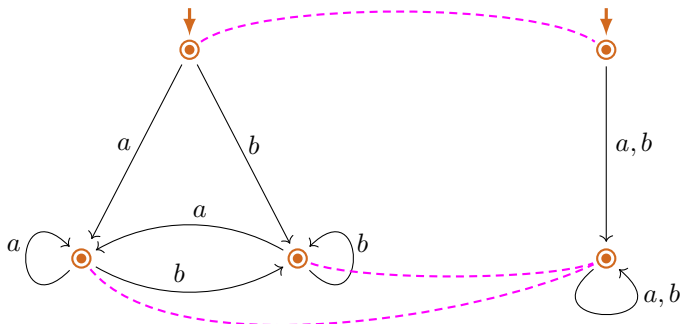


**LEE**



**LEE**

# Chart interpretation (example)

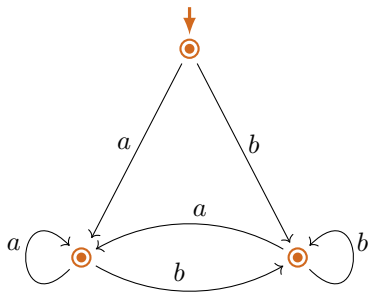


$$\mathcal{C}((a^* \cdot b^*)^*) \quad \leftrightarrow \quad \mathcal{C}((a + b)^*)$$

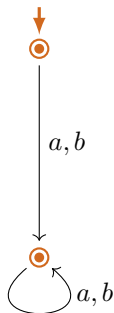
$$\llbracket (a^* \cdot b^*)^* \rrbracket_{\mathcal{P}} \quad = \quad \llbracket (a + b)^* \rrbracket_{\mathcal{P}}$$



# Chart interpretation (example)

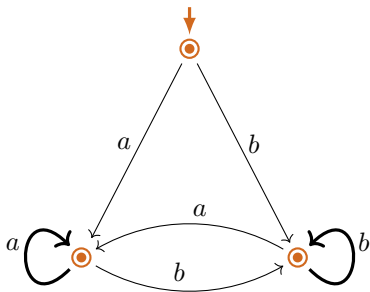


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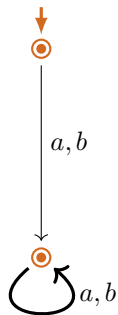


$$\mathcal{C}((a + b)^*)$$

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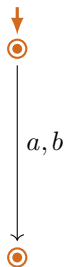
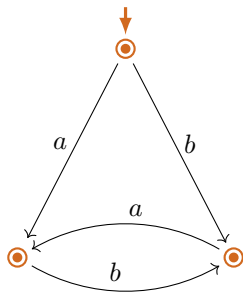


$$\mathcal{C}((a^* \cdot b^*)^*)$$



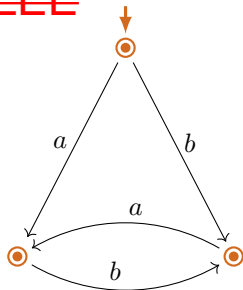
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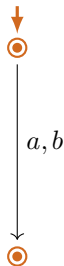


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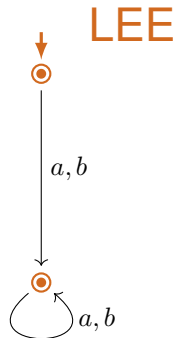
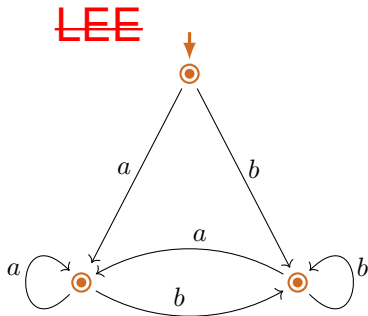
LEE



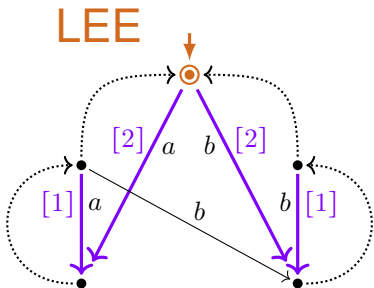
LEE



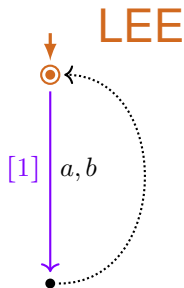
# Chart interpretation (example)



# LEE-1-chart interpretation (example) (TERMGRAPH, 2020)

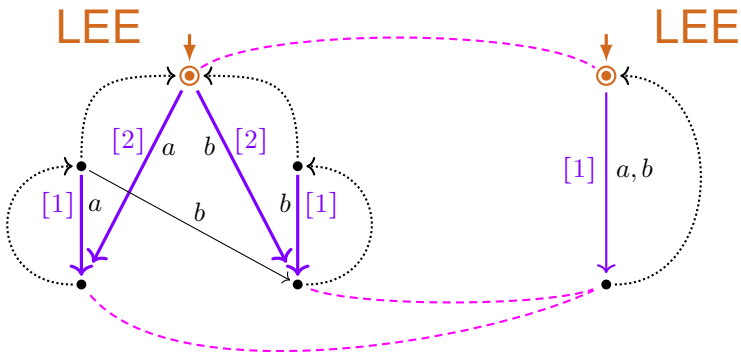


$$\underline{C}((a^* \cdot b^*)^*)$$



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# LEE-1-chart interpretation (example) (TERMGRAPH, 2020)

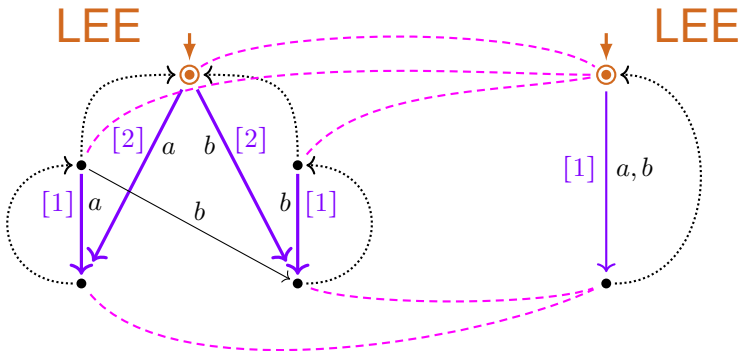


$$\underline{\mathcal{C}}((a^* \cdot b^*)^*)$$

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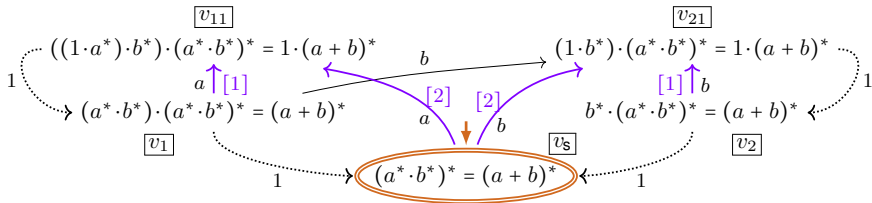
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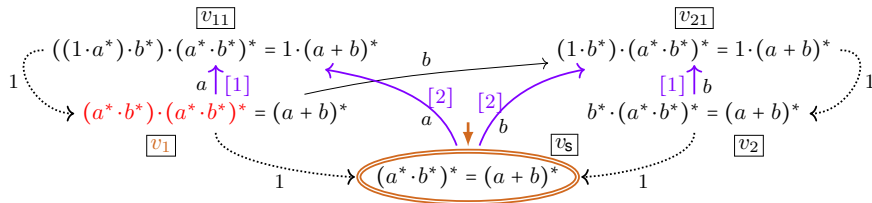


# LEE-witnessed coinductive proof over Mil<sup>-</sup>



Right- and left-hand sides are Mil<sup>-</sup>-provable solutions in every vertex.

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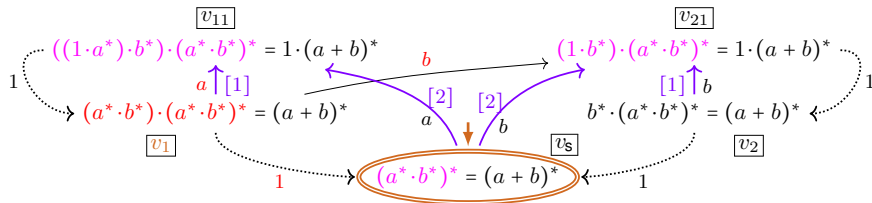


Right- and left-hand sides are Mil<sup>-</sup>-provable solutions in every vertex.

E.g. in  $v_1$ :

$$(a^* \cdot b^*) \cdot (a^* \cdot b^*)^* =_{\text{Mil}^-}$$

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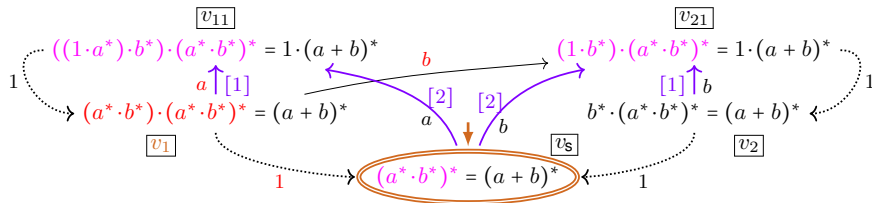


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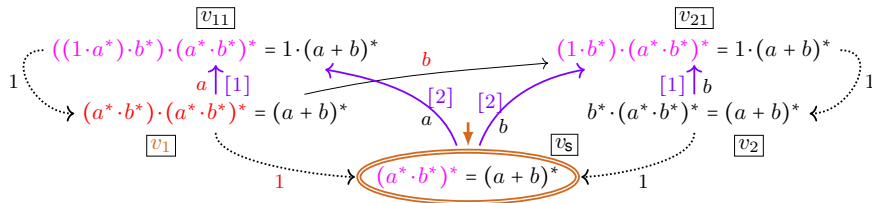
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$$(a^* \cdot b^*) \cdot (a^* \cdot b^*)^* =_{\text{Mil}^-}$$

$$=_{\text{Mil}^-} 1 \cdot (a^* \cdot b^*)^* + a \cdot (((1 \cdot a^*) \cdot b^*) \cdot (a^* \cdot b^*)^*) + b \cdot ((1 \cdot b^*) \cdot (a^* \cdot b^*)^*)$$

# LEE-witnessed coinductive proof over Mil<sup>-</sup>

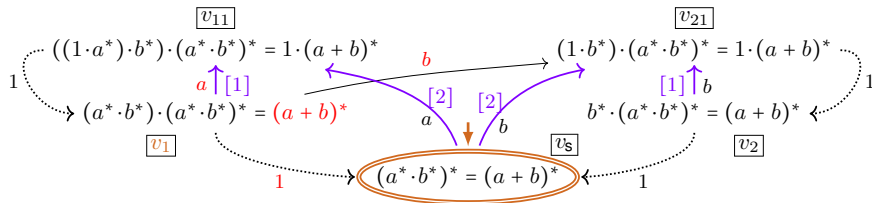


Right- and left-hand sides are **Mil<sup>-</sup>-provable solutions** in every vertex.

E.g. in  $v_1$ :

$$\begin{aligned}
 (a^* \cdot b^*) \cdot (a^* \cdot b^*)^* &=_{\text{Mil}^-} ((1 + a \cdot a^*) \cdot (1 + b \cdot b^*)) \cdot (a^* \cdot b^*)^* \\
 &=_{\text{Mil}^-} (1 \cdot 1 + a \cdot a^* \cdot 1 + 1 \cdot b \cdot b^* + a \cdot a^* \cdot b \cdot b^*) \cdot (a^* \cdot b^*)^* \\
 &=_{\text{Mil}^-} (1 + a \cdot a^* + a \cdot a^* \cdot b \cdot b^* + b \cdot b^*) \cdot (a^* \cdot b^*)^* \\
 &=_{\text{Mil}^-} (1 + a \cdot a^* \cdot (1 + b \cdot b^*) + b \cdot b^*) \cdot (a^* \cdot b^*)^* \\
 &=_{\text{Mil}^-} (1 + a \cdot a^* \cdot b^* + b \cdot b^*) \cdot (a^* \cdot b^*)^* \\
 &=_{\text{Mil}^-} 1 \cdot (a^* \cdot b^*)^* + a \cdot (((1 \cdot a^*) \cdot b^*) \cdot (a^* \cdot b^*)^*) + b \cdot ((1 \cdot b^*) \cdot (a^* \cdot b^*)^*)
 \end{aligned}$$

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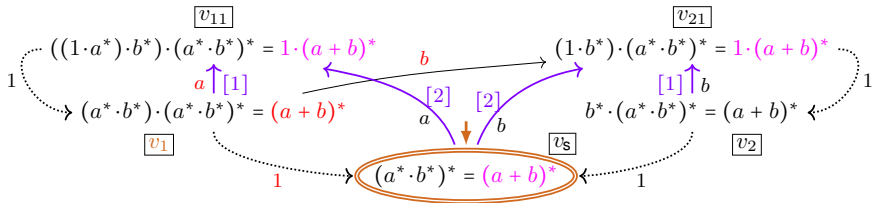


Right- and left-hand sides are Mil<sup>-</sup>-provable solutions in every vertex.

E.g. in  $v_1$ :

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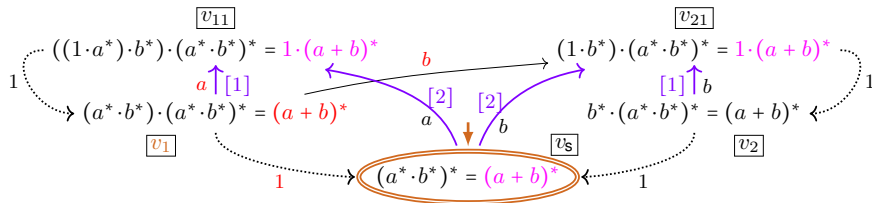


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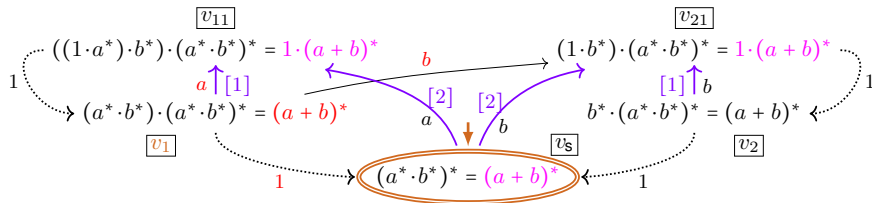
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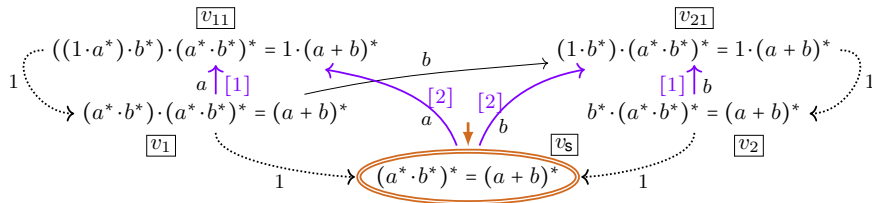


Right- and left-hand sides are Mil<sup>-</sup>-provable solutions in every vertex.

E.g. in  $v_1$ :

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 (a+b)^* &=_{\text{Mil}^-} (a+b)^* + (a+b)^* =_{\text{Mil}^-} 1 + (a+b) \cdot (a+b)^* + 1 + (a+b) \cdot (a+b)^* \\
 &=_{\text{Mil}^-} 1 + 1 + (a+b) \cdot (a+b)^* + a \cdot (a+b)^* + b \cdot (a+b)^* \\
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 &=_{\text{Mil}^-} 1 \cdot (a+b)^* + a \cdot (1 \cdot (a+b)^*) + b \cdot (1 \cdot (a+b)^*)
 \end{aligned}$$

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Right- and left-hand sides are Mil<sup>-</sup>-provable solutions in every vertex.

# Coinductive proof systems

Rule scheme for combining LEE-witnessed coinductive proofs:

$$\frac{g_1 = h_1 \quad \dots \quad g_n = h_n \quad \mathcal{LCP}(e = f)}{e = f} \text{LCoProof}_n$$

- ▶  $\mathcal{LCP}(e = f)$  is a LEE-witnessed coinductive proof of  $e = f$  over  $\text{Mil}^- + \{g_1 = h_1, \dots, g_n = h_n\}$ .

We define the proof systems:

$\text{CLC} := \text{rules } \{\text{LCoProof}_n\}_{n \in \mathbb{N}}$

$\text{cMil} := \text{Mil}^- + \{\text{LCoProof}_n\}_{n \in \mathbb{N}}$

Lemma

$\text{CLC} \sim \text{cMil}$

# Coinductive proof systems

Rule scheme for **combining coinductive proofs**:

$$\frac{g_1 = h_1 \quad \dots \quad g_n = h_n \quad \mathcal{CP}(e = f)}{e = f} \text{CoProof}_n$$

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over  $\text{Mil}^- + \{g_1 = h_1, \dots, g_n = h_n\}$ .

We define the proof systems:

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$\text{CC} := \text{rules } \{\text{CoProof}_n\}_{n \in \mathbb{N}}$

$\text{cMil} := \text{Mil}^- + \{\text{LCoProof}_n\}_{n \in \mathbb{N}}$

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Lemma

$\text{CLC} \sim \text{cMil}$

Lemma

- (i)  $\text{CC} \sim \overline{\text{cMil}}$ .
- (ii)  $\overline{\text{cMil}} \sim \text{Mil}^- + \text{USP}$ ,
- (iii)  $\text{CC}$  and  $\overline{\text{cMil}}$   
are complete for  $=_{\llbracket \cdot \rrbracket_P}$ .

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Lemma

$\text{CLC} \sim \text{cMil}$

Consequence

$\text{cMil} \sim \text{CLC} \approx \text{CC} \sim \overline{\text{cMil}}$ .

Lemma

- (i)  $\text{CC} \sim \overline{\text{cMil}}$ .
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# Proof transformation Mil $\mapsto$ cMil

$$\begin{array}{c}
 \cdot \\
 \frac{e = f \cdot e + g}{e = f^* \cdot g} \text{RSP}^* \\
 \Longrightarrow \frac{e = f \cdot e + g \quad \mathcal{LCP}_{\text{Mil}^- + \{e = f \cdot e + g\}}(e = f^* \cdot g)}{e = f^* \cdot g} \text{LCoProof}_1
 \end{array}$$

# Proof transformation Mil $\mapsto$ cMil

$$\frac{\overbrace{(a+b)^*}^e = \overbrace{((a \cdot a^* + b) \cdot b^*)}^f \cdot \overbrace{(a+b)^*}^e + \overbrace{1}^g}}{\overbrace{(a+b)^*}^e = \underbrace{((a \cdot a^* + b) \cdot b^*)^*}_{f^*} \cdot \underbrace{1}_g} \text{RSP}^*$$

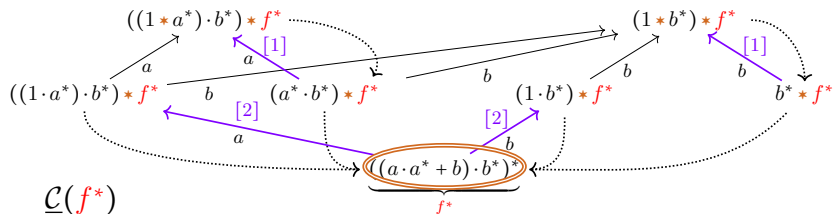
# From fixed-point rule instances to coinductive proofs

$$\frac{\overbrace{(a+b)^*}^e = \overbrace{((a \cdot a^* + b) \cdot b^*)}^f \cdot \overbrace{(a+b)^*}^e + \overbrace{1}^g}}{\overbrace{(a+b)^*}^e = \underbrace{((a \cdot a^* + b) \cdot b^*)^*}_{f^*} \cdot \underbrace{1}_g} \text{RSP}^*$$



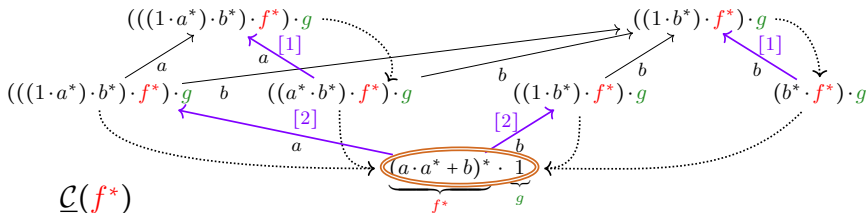
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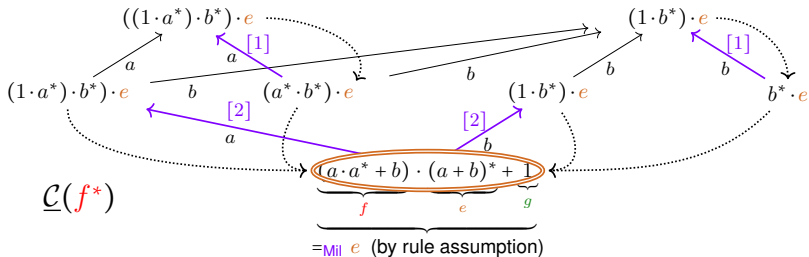
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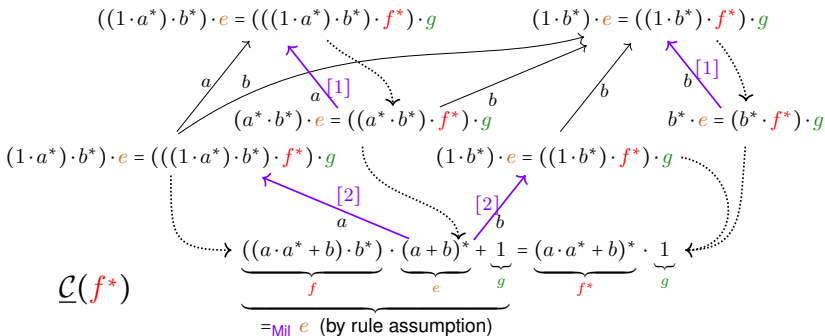
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# From fixed-point rule instances to coinductive proofs

$$\frac{\overbrace{(a+b)^*}^e = \overbrace{((a \cdot a^* + b) \cdot b^*)}^f \cdot \overbrace{(a+b)^*}^e + \overbrace{1}^g}}{\overbrace{(a+b)^*}^e = \underbrace{((a \cdot a^* + b) \cdot b^*)^*}_{f^*} \cdot \underbrace{1}_g}} \text{RSP}^*$$



LEE-witnessed coinductive proof over  $\text{Mil}^- + \{e = f \cdot e + g\}$

# Proof transformation $\text{Mil} \mapsto \text{cMil}$

$$\begin{array}{c}
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 \frac{e = f \cdot e + g}{e = f^* \cdot g} \text{RSP}^* \\
 \Longrightarrow \frac{e = f \cdot e + g \quad \mathcal{LCP}_{\text{Mil}^- + \{e = f \cdot e + g\}}(e = f^* \cdot g)}{e = f^* \cdot g} \text{LCoProof}_1
 \end{array}$$

# Proof transformation $\text{Mil} \mapsto \text{cMil}$

## Theorem

$\text{Mil} \approx \text{cMil}$ , because:

every derivation in  $\text{Mil}$  with conclusion  $e = f$  can be transformed effectively into a derivation in  $\text{cMil}$  with conclusion  $e = f$ .

## Proof idea.

$$\frac{e = f \cdot e + g}{e = f^* \cdot g} \text{RSP}^* \quad \Longrightarrow \quad \frac{e = f \cdot e + g \quad \mathcal{LCP}_{\text{Mil}^- + \{e = f \cdot e + g\}}(e = f^* \cdot g)}{e = f^* \cdot g} \text{LCoProof}_1$$

## Corollary

$\text{Mil} \sim \text{cMil} \sim \text{CLC}$ .

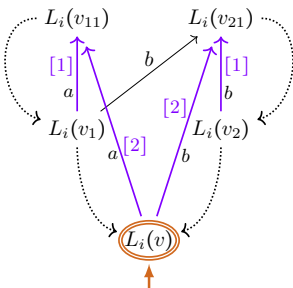
# Proof transformation $\text{cMil} \mapsto \text{Mil}$

## Lemma

For all star expression  $e, f$ , and equations  $\Gamma \subseteq =_{\text{Mil}}$  :

$$e \stackrel{\text{LEE}}{=}_{\text{Mil} + \Gamma} f \quad \Longrightarrow \quad e =_{\text{Mil}} f$$

# Extraction of Mil-derivation from LEE-witn. coind. proof

 $\underline{C}, \hat{C}$ 


$$L_i(v_{21}) = \stackrel{(\text{sol})}{\text{Mil}^-} 1 \cdot L_i(v_2) =_{\text{Mil}^-} L_i(v_2)$$

( $\stackrel{(\text{sol})}{\text{Mil}^-}$  means use of 'is  $\text{Mil}^-$ -provable solution')

$$L_i(v_2) = \stackrel{(\text{sol})}{\text{Mil}^-} b \cdot L_i(v_{21}) + 1 \cdot L_i(v) =_{\text{Mil}^-} b \cdot L_i(v_2) + L_i(v)$$

$\Downarrow$  applying **RSP\***

$$L_i(v_2) =_{\text{Mil}} b^* \cdot L_i(v)$$

$$L_i(v_{11}) = \stackrel{(\text{sol})}{\text{Mil}^-} 1 \cdot L_i(v_1) =_{\text{Mil}^-} L_i(v_1)$$

$$\begin{aligned} L_i(v_1) &=_{\text{Mil}^-} a \cdot L_i(v_{11}) + b \cdot L_i(v_{21}) + 1 \cdot L_i(v) \\ &=_{\text{Mil}} a \cdot L_i(v_1) + (b \cdot b^* + 1) \cdot L_i(v) \\ &=_{\text{Mil}^-} a \cdot L_i(v_1) + b^* \cdot L_i(v) \end{aligned}$$

$\Downarrow$  applying **RSP\***

$$L_i(v_1) =_{\text{Mil}} a^* \cdot (b^* \cdot L_i(v)) =_{\text{Mil}^-} (a^* \cdot b^*) \cdot L_i(v)$$

$$\begin{aligned} L_i(v) &= \stackrel{(\text{sol})}{\text{Mil}^-} 1 + a \cdot L_i(v_{11}) + b \cdot L_i(v_{21}) =_{\text{Mil}^-} 1 + a \cdot L_i(v_1) + b \cdot L_i(v_2) \\ &=_{\text{Mil}} (a \cdot (a^* \cdot b^*) + b \cdot b^*) \cdot L_i(v) + 1 \end{aligned}$$

$\Downarrow$  applying **RSP\***

$$L_i(v) =_{\text{Mil}} (a \cdot (a^* \cdot b^*) + b \cdot b^*)^* \cdot 1 =_{\text{Mil}^-} (a \cdot (a^* \cdot b^*) + b \cdot b^*)^* = \underline{s_{\hat{C}}}(v)$$



# Proof transformation $\text{cMil} \mapsto \text{Mil}$

## Lemma (extraction and unique solvability)

Let  $\underline{C}$  be a LEE-1-chart.

- ▶ From  $\underline{C}$  a  $\text{Mil}^-$ - (hence  $\text{Mil}$ -) provable solution can be extracted.
- ▶ Any two  $\text{Mil}$ -provable solutions of  $\underline{C}$  are  $\text{Mil}$ -provably equal.

## Lemma

For all star expression  $e, f$ , and equations  $\Gamma \subseteq =_{\text{Mil}}$  :

$$e \stackrel{\text{LEE}}{=}_{\text{Mil}^+ + \Gamma} f \quad \Longrightarrow \quad e =_{\text{Mil}} f$$

## Theorem

$\text{cMil} \lesssim \text{Mil}$ , because:

every derivation in  $\text{cMil}$  with conclusion  $e = f$  can be transformed effectively into a derivation in  $\text{Mil}$  with conclusion  $e = f$ .

# Summary

We define:

- ▶ (LEE-witnessed) coinductive proofs over  $Mil^-$  :
  - ▶ 1-charts  $\underline{C}$  (with LEE) whose vertices are labeled by equations between the values of two provable solutions of  $\underline{C}$
- ▶ proof systems
  - ▶ systems cMil / CLC with LEE-witnessed coind. proofs over  $Mil^-$
  - ▶ systems  $\overline{cMil}$  / CC with coinductive proofs over  $Mil^-$

Results:

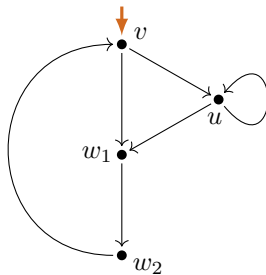
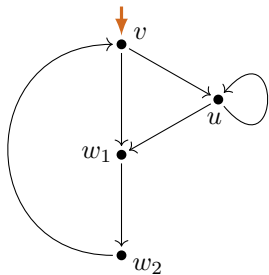
- ▶  $Mil = (Mil^- + RSP^*) \sim (Mil^- + LCoProof) = cMil \sim CLC$
- ▶  $Mil \approx (Mil^- + USP) \sim (Mil^- + CoProof) = \overline{cMil} \sim CC$  ((clearly) complete).

Desired application: proof strategy for completeness proof of Mil

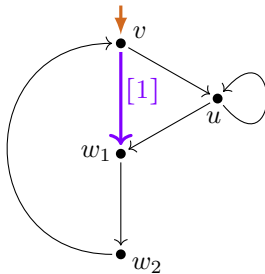
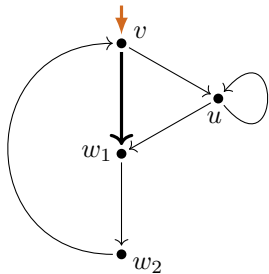
$$\text{▶ } \vdash_{Mil} e = f \iff \vdash_{cMil} e = f \iff e = \llbracket \cdot \rrbracket_P f$$

- 
- ▶ **Technical report:** [arXiv:2108.13104](https://arxiv.org/abs/2108.13104)

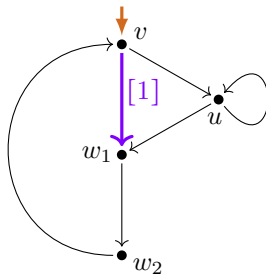
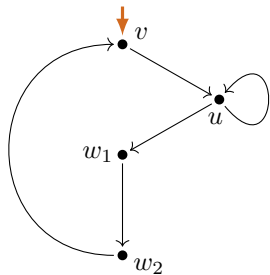
# Layered LEE-witness



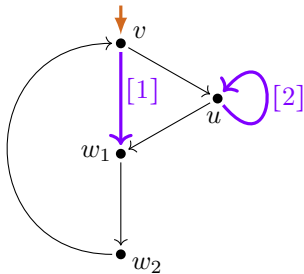
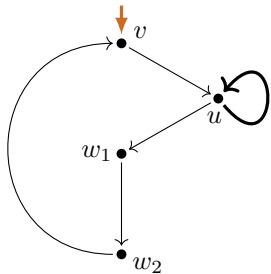
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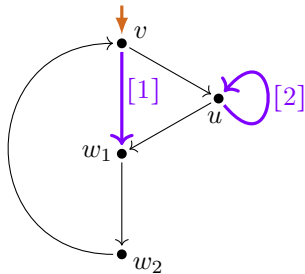
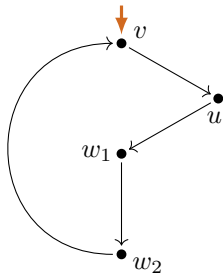
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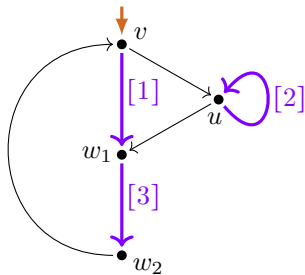
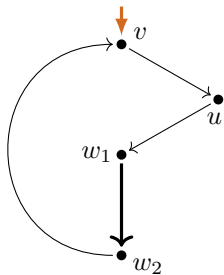
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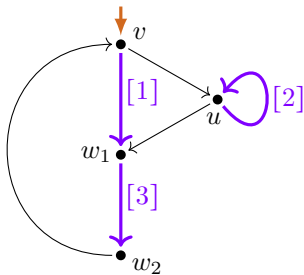
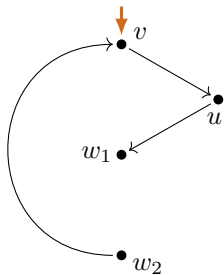


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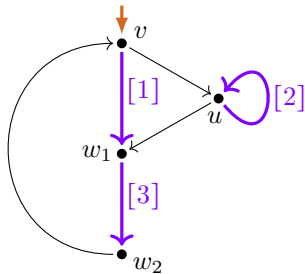
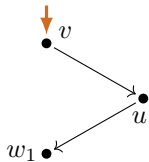




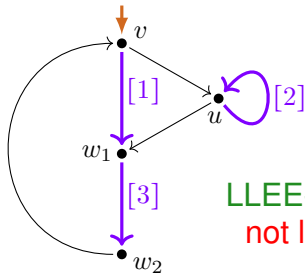
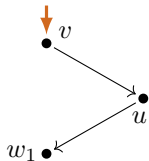
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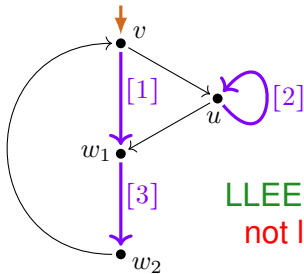
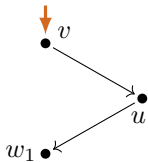


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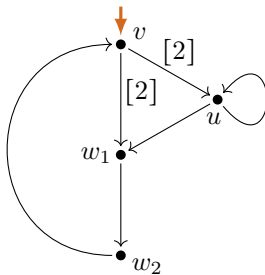
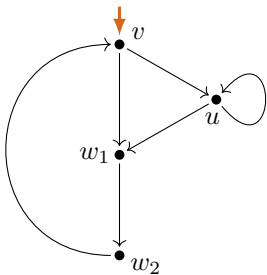


LEE-witness  
not layered

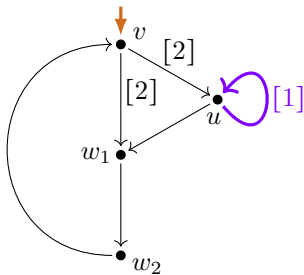
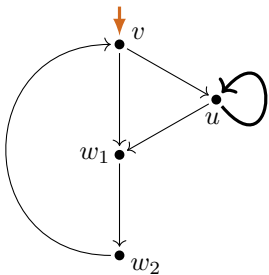
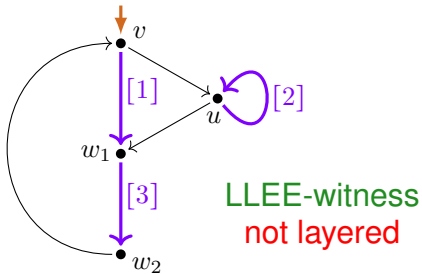
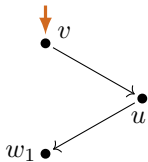
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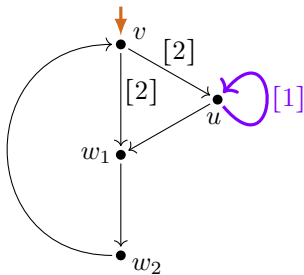
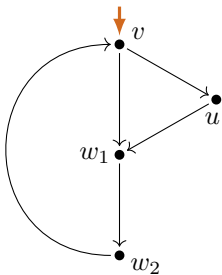
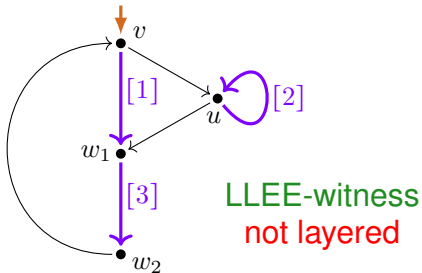
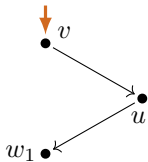
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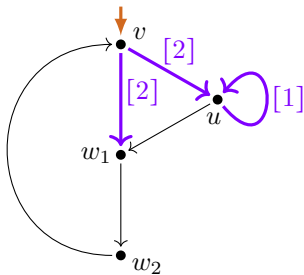
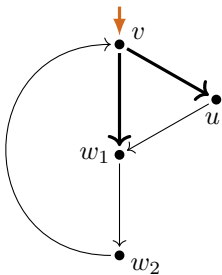
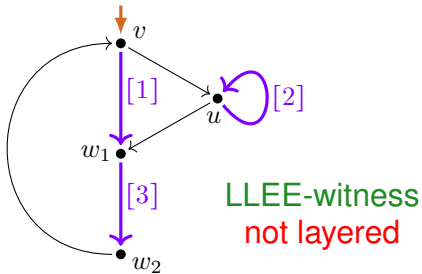
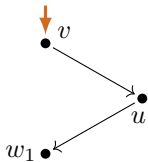
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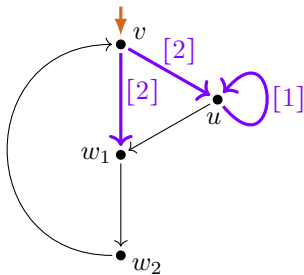
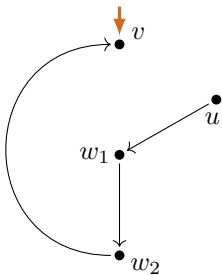
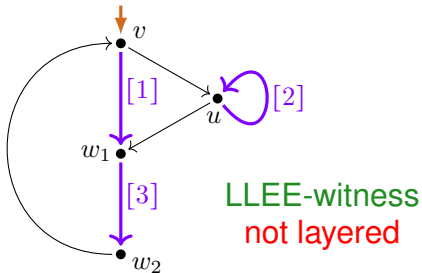
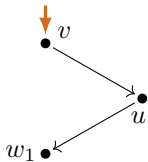
# Layered LEE-witness



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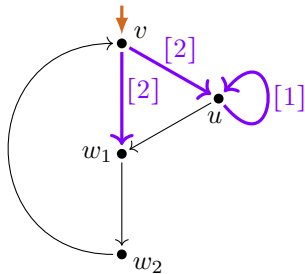
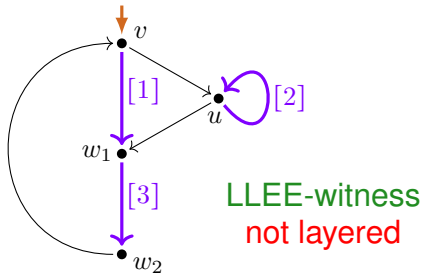
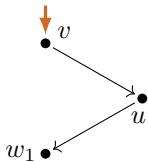


# Layered LEE-witness

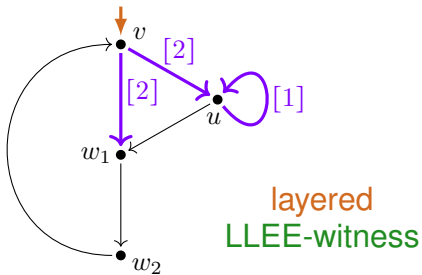
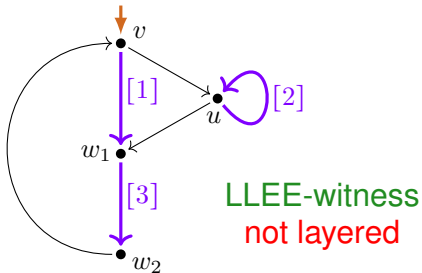
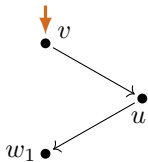




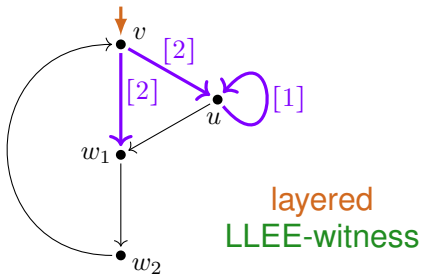
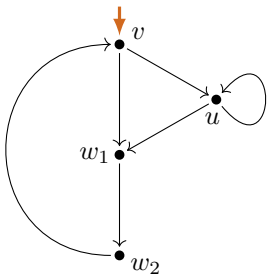
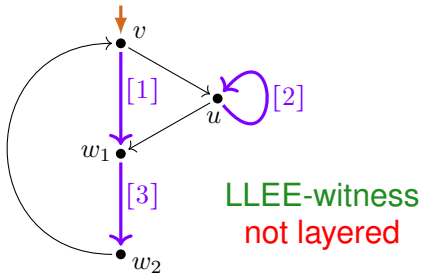
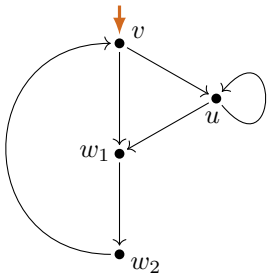
# Layered LEE-witness



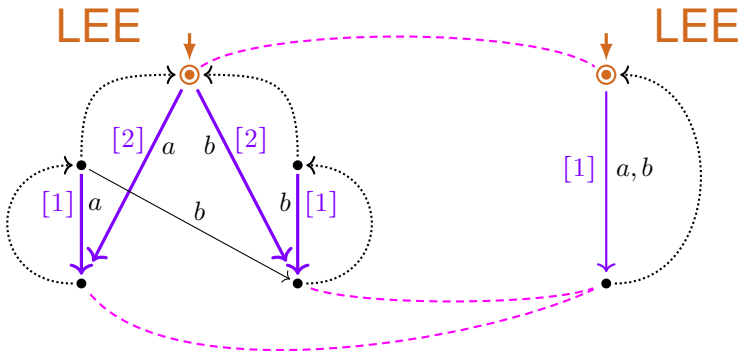
# Layered LEE-witness



# Layered LEE-witness



# LEE-1-chart interpretation (example) (TERMGRAPH, 2020)

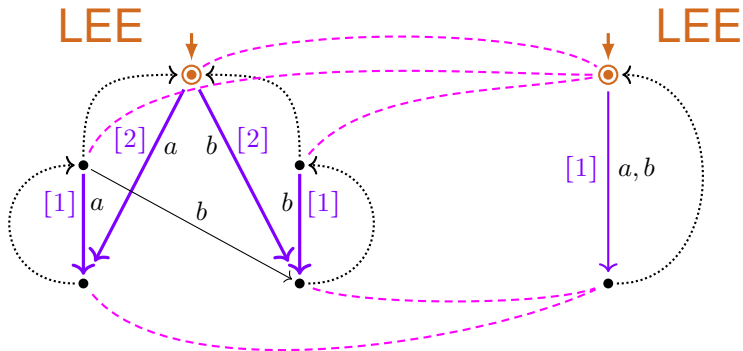


$$\underline{\mathcal{C}}((a^* \cdot b^*)^*)$$

$$\longleftrightarrow^{(1)}$$

$$\underline{\mathcal{C}}((a + b)^*)$$

# LEE-1-chart interpretation (example) (TERMGRAPH, 2020)

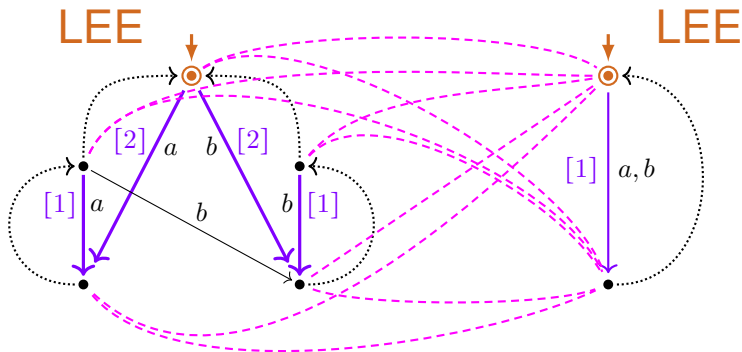


$$\underline{\mathcal{C}}((a^* \cdot b^*)^*)$$

$$\underline{\leftrightarrow}^{(1)}$$

$$\underline{\mathcal{C}}((a + b)^*)$$

# LEE-1-chart interpretation (example) (TERMGRAPH, 2020)

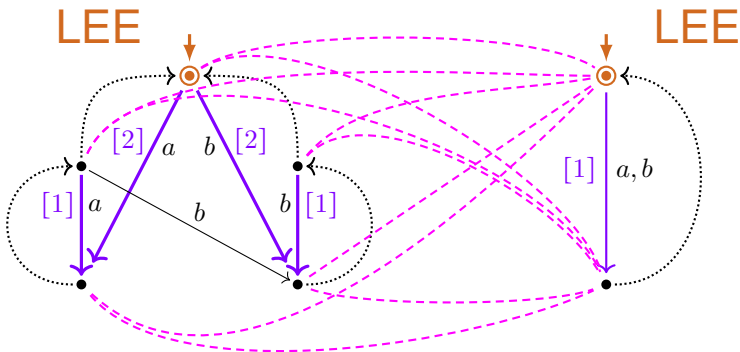


$$\underline{\mathcal{C}}((a^* \cdot b^*)^*)$$

$$\longleftrightarrow^{(1)}$$

$$\underline{\mathcal{C}}((a + b)^*)$$

# LEE-1-chart interpretation (example) (TERMGRAPH, 2020)

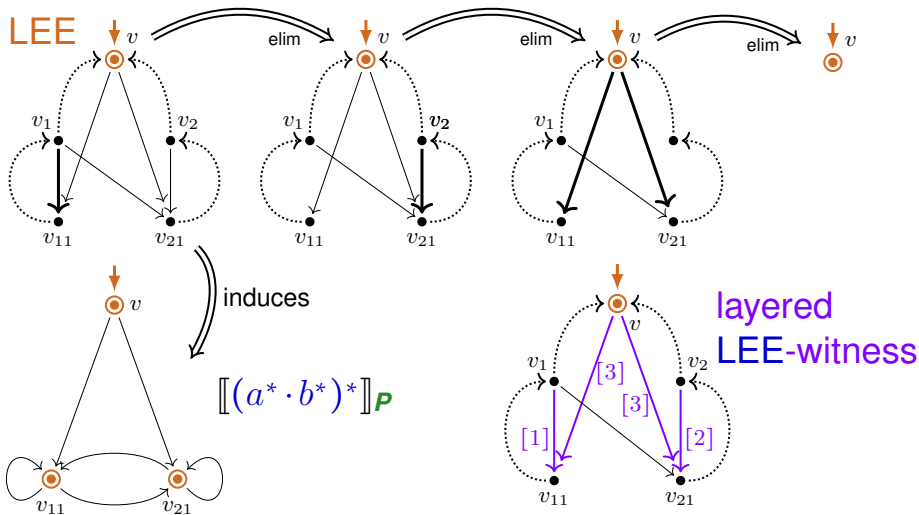


$$\underline{\mathcal{C}}((a^* \cdot b^*)^*)$$

$$\overset{(1)}{\longleftrightarrow}$$

$$\underline{\mathcal{C}}((a + b)^*)$$

# LEE, and LLEE-witness, induced process graph





# LEE-charts: properties and results

## Lemmas

- (I) Chart interpretations of **1-free star expressions** satisfy **LEE**.
- (SU) **LEE**-charts have unique provable solutions up to **Mil**-provability.
- (C) **LEE** is preserved under bisimulation collapse.

## Theorem (G/Fokkink, LICS 2020)

The adaptation **BBP** of **Mil** to **1-free star expressions** is complete.

## Consequence of lemmas used

- (E) A chart  $\mathcal{C}$  is **expressible by a 1-free star expr. modulo bisimilarity**  
 $\iff$  the bisimulation collapse of  $\mathcal{C}$  is a **LEE**-chart.

Hence expressible | not expressible by **1-free star expressions**:

