

Optics, functorially

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- **Monads**: model notions of computations
[Moggi'89, '91]
- **Monad transformers**: combine computational effects
[Liang-Hudak-Jones'95, Moggi'97, Benton-Hughes-Moggi'00]
- **Lenses**: bidirectional (bx) transformations (**view-update** of databases)
($\text{get} : A \rightarrow B, \text{put} : A \times B' \rightarrow A'$) [Oles'82]
- **Lens transformers?**
- More generally, **(mixed) optics** as bx transformations

Features: modularity, compositionality

Applications: game theory, machine learning, database systems

Therefore ... what about **optics transformers?**

Name	Description	Actions	Base
Adapter	$\mathbf{C}(S, A) \otimes \mathbf{D}(B, T)$	(Optic _{id, id})	\mathcal{V}, \otimes
Lens	$\mathbf{C}(S, A) \times \mathbf{D}(S \bullet B, T)$	(Optic _{\times, \bullet})	\mathcal{W}, \times
Monoidal lens	$\mathbf{CCom}(S, A) \times \mathbf{C}(US \otimes B, T)$	(Optic _{\otimes, \times})	\mathcal{W}, \times
Algebraic lens	$\mathbf{C}(S, A) \times \mathbf{D}(\Psi S \bullet B, T)$	(Optic _{$\iota \times, \iota \bullet$})	\mathcal{W}, \times
Monadic lens	$\mathcal{W}(S, A) \times \mathcal{W}(S \times B, \Psi T)$	(Optic _{\times, \times})	\mathcal{W}, \times
Linear lens	$\mathbf{C}(S, [B, T] \bullet A)$	(Optic _{\bullet, \otimes})	\mathcal{V}, \otimes
Prism	$\mathbf{C}(S, T \bullet A) \times \mathbf{D}(B, T)$	(Optic _{$\bullet, +$})	\mathcal{W}, \times
Coalg. prism	$\mathbf{C}(S, \Theta T \bullet A) \times \mathbf{D}(B, T)$	(Optic _{$\iota \bullet, \iota +$})	\mathcal{W}, \times
Grate	$\mathbf{D}([S, A] \bullet B, T)$	(Optic _{$\{ \cdot \}, \bullet$})	\mathcal{V}, \otimes
Glass	$\mathbf{C}(S \times [[S, A], B], T)$	(Optic _{$\times [\cdot], \times [\cdot], \cdot$})	\mathcal{W}, \times
Affine traversal	$\mathbf{C}(S, T + A \otimes \{B, T\})$	(Optic _{$+ \otimes, + \otimes$})	\mathcal{W}, \times
Traversal	$\mathcal{V}(S, \sum^n A^n \otimes [B^n, T])$	(Optic _{Pw, Pw})	\mathcal{V}, \otimes
Kaleidoscope	$\sum^n \mathcal{V}([A^n, B], [S^n, T])$	(Optic _{App, App})	\mathcal{V}, \otimes
Setter	$\mathcal{V}([A, B], [S, T])$	(Optic _{ev, ev})	\mathcal{V}, \otimes
Fold	$\mathcal{V}(S, \mathcal{L}A)$	(Optic _{Foldable, \bullet})	\mathcal{V}, \otimes

Table of optics [Clarke et al. '24]

Monoidal category: a category \mathcal{M} equipped with a tensor product (bifunctor)

$$\mathcal{M} \times \mathcal{M} \longrightarrow \mathcal{M}, \quad (M, N) \mapsto M \otimes N$$

and a unit object I , such that \otimes is associative and unital up to coherent isomorphism

Examples

- Any category \mathcal{M} with finite (co)products
- Endofunctors $[\mathcal{A}, \mathcal{A}]$, with functor composition
- Presheaves $[\mathcal{A}^{\text{op}}, \text{Set}]$ over a (small) category \mathcal{A} , with Day convolution
- The Eilenberg-Moore category of algebras of a commutative monad on a monoidal category

[Beñabou'67, McCrudden'00, Capucci-Gavranović'22]

$(\mathcal{M}, \otimes, I)$ monoidal category

\mathcal{M} -actegory: a category \mathcal{A} equipped with an action (bifunctor)

$$\mathcal{M} \times \mathcal{A} \longrightarrow \mathcal{A}, (M, A) \mapsto M \cdot A$$

associative and unital up to coherent isomorphism

$$(M \otimes N) \cdot A \cong M \cdot (N \cdot A), I \cdot A \cong A$$

Examples

- Any monoidal category \mathcal{M} acts on itself via the tensor product
- For any category \mathcal{A} , $[\mathcal{A}, \mathcal{A}]$ acts on \mathcal{A} via functor application
- The Kleisli category of a strong monad on a monoidal category

Lax \mathcal{M} -morphism: functor $F : \mathcal{A} \rightarrow \mathcal{B}$ between \mathcal{M} -actegories, endowed with natural transformation $st : M \cdot FA \rightarrow F(M \cdot A)$ (**strength**), compatible with the \mathcal{M} -actions

Example. If \mathcal{M} is a cartesian category acting on itself, then an \mathcal{M} -lax endofunctor $\mathcal{M} \rightarrow \mathcal{M}$ is precisely a **strong** functor.

Remark. A **lax \mathcal{M} -morphism structure** on an endofunctor F on an \mathcal{M} -actegory is the same as a **lifting** of the \mathcal{M} -action to $\text{Coalg}(F)$, such that the forgetful functor becomes strict \mathcal{M} -morphism (strength is identity)

Colax \mathcal{M} -morphism: lax \mathcal{M} -morphism between opposite actegories. A colax \mathcal{M} -morphism F comes equipped with a **costrength** $cst : F(M \cdot A) \rightarrow M \cdot FA$

Example. Let \mathcal{M} be a cartesian category acting on itself. An endofunctor on \mathcal{M} is a colax \mathcal{M} -morphism (also known as **costrong** functor) iff it is **copointed** [B-Pantelimon'24]

(Co)Lax \mathcal{M} -transformation: natural transformation between (co)lax \mathcal{M} -morphisms, compatible with their (co)strengths

The Para construction

[Wood'76, Hermida-Tennent'12, Fong-Spivak-Tuyéras'19, Capucci-Gavranović-Hedges'20]

Let \mathcal{A} be an \mathcal{M} -actegory

Para(\mathcal{A}): bicategory which “adds parameters” to \mathcal{A}

- the objects are those of \mathcal{A}
- the morphisms are M -parametrised morphisms ($M \in \mathcal{M}, f : M \cdot A \longrightarrow B$)
- 2-cells are given by **reparametrisation**

Remarks

- The construction $\mathcal{M}\text{-Act}_l \longrightarrow \mathbf{Bicat}, \mathcal{A} \mapsto \mathbf{Para}(\mathcal{A})$ is functorial with respect to **lax \mathcal{M} -morphisms**.
- There is a 2-opfibration $\mathbf{Para}(\mathcal{A}) \longrightarrow \mathbf{B.M}$ over the delooping of \mathcal{M} , by projecting parameters.
- Dually, **CoPara**(\mathcal{A}) = **Para**(\mathcal{A}^{op}) gives a functor $\mathcal{M}\text{-Act}_c \longrightarrow \mathbf{Bicat}$ and a 2-fibration over $\mathbf{B.M}$

Informally: optics are coupled pairs of coparametrised and parametrised morphisms, but with externally **unobservable** joint parameter

Let \mathcal{A}, \mathcal{B} be two \mathcal{M} -actegories

A **(mixed) optic** from (A, B) with the **focus** on (A', B') is an element of the coend

$$\mathbf{Optic}_{\mathcal{A}, \mathcal{B}}((A, B), (A', B')) = \int^{\mathcal{M}} \mathcal{A}(A, M \cdot A') \times \mathcal{B}(M \cdot B', B)$$

where $A, A' \in \mathcal{A}$ and $B, B' \in \mathcal{B}$

Optics are arrows of a **category** $\mathbf{Optic}_{\mathcal{A}, \mathcal{B}}$, which comes with an identity on objects fully faithful functor

$$\mathcal{A}^{\text{op}} \times \mathcal{B} \longrightarrow \mathbf{Optic}_{\mathcal{A}, \mathcal{B}}$$

Lenses arise when $\mathcal{M} = \mathcal{A} = \mathcal{B}$ and both the monoidal structure and the actions are given by the cartesian product

$$\int^{\mathcal{M}} \mathcal{M}(A, M \times A') \times \mathcal{M}(M \times B', B) \cong \mathcal{M}(A, A') \times \mathcal{M}(A \times B', B)$$

Optic $_{\mathcal{A}, \mathcal{B}}$ as a **bicategory**:
[Braithwaite et al.'21]

$$\begin{array}{ccc}
 \mathbf{Optic}_{\mathcal{A}, \mathcal{B}} & \longrightarrow & \mathbf{Para}(\mathcal{B}) \\
 \downarrow & \lrcorner & \downarrow \\
 \mathbf{CoPara}(\mathcal{A}) & \longrightarrow & \mathbf{B}\mathcal{M}
 \end{array}$$

The 1-cells in the bicategory **Optic** $_{\mathcal{A}, \mathcal{B}}$ are pairs of coparametrised, respectively parametrised morphisms

$$(M, f : A \longrightarrow M \cdot A', g : M \cdot B' \longrightarrow B)$$

explicitly keeping track of the residual, without constraints (which now live in the 2-cells of **Optic** $_{\mathcal{A}, \mathcal{B}}$)

The functors **Para** : $\mathcal{M}\text{-Act}_I \longrightarrow \mathbf{Bicat}$ and **CoPara** : $\mathcal{M}\text{-Act}_c \longrightarrow \mathbf{Bicat}$ consequently induce a pseudofunctor [B-Pantelimon'24]

$$\mathbf{Optic} : \mathcal{M}\text{-Act}_c \otimes \mathcal{M}\text{-Act}_I \longrightarrow \mathbf{Bicat}$$

Let $\mathcal{M} = \mathcal{A} = \mathcal{B}$, with monoidal structure and \mathcal{M} -actions given by the cartesian product

There is a **local adjunction** between the bicategory of optics and the (discrete) bicategory of lenses [Gavranović '22]

$$\mathbf{Lens}_{\mathcal{M}}((A, A'), (B, B')) \begin{array}{c} \xrightarrow{\quad \perp \quad} \\ \xleftarrow{\quad \perp \quad} \end{array} \mathbf{Optic}_{\mathcal{M}, \mathcal{M}}((A, A'), (B, B'))$$

which becomes a **bijection** of homsets when restricted to the 1-category of optics.

Let $F, G : \mathcal{M} \rightarrow \mathcal{M}$ be costrong, respectively strong functors

Then there is a **morphism of adjunctions** given by $\mathbf{Lens}(F, G)$ and $\mathbf{Optic}(F, G)$ [B-Pantelimon'24]

$$\begin{array}{ccc} \mathbf{Lens}_{\mathcal{M}}((A, A'), (B, B')) & \begin{array}{c} \xrightarrow{\quad \perp \quad} \\ \xleftarrow{\quad \perp \quad} \end{array} & \mathbf{Optic}_{\mathcal{M}, \mathcal{M}}((A, A'), (B, B')) \\ \mathbf{Lens}(F, G) \downarrow & & \downarrow \mathbf{Optic}(F, G) \\ \mathbf{Lens}_{\mathcal{M}}((FA, GA'), (FB, GB')) & \begin{array}{c} \xrightarrow{\quad \perp \quad} \\ \xleftarrow{\quad \perp \quad} \end{array} & \mathbf{Optic}_{\mathcal{M}, \mathcal{M}}((FA, GA'), (FB, GB')) \end{array}$$

Dependent optics [Vertechni'22] generalise (mixed) optics, replacing the actions of the monoidal category \mathcal{M} by a pair of pseudofunctors

$$\mathbf{B}^{\text{op}} \longrightarrow \mathbf{Cat}$$

where \mathbf{B} is a bicategory (taking \mathbf{B} to be the delooping of \mathcal{M} recovers usual optics)

Our previous **Optic** construction **extends to dependent optics** when lax and colax \mathcal{M} -morphisms are replaced by lax and colax natural transformations.

- Better understand optics properties deriving from their inherent fibrational nature
- Instantiate **Optic**(F, G) to various classes of optics (prisms, traversables, etc.); gain more intuition on these
- Find (if any) connections with monad-comonad interaction laws

Thank you for your attention!