

Graded Semantics and Graded Logics for Eilenberg-Moore Coalgebras

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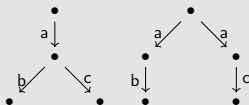
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Context

Behavioural Equivalence

Often too fine grained

$$\begin{array}{ccccc}
 X & \longrightarrow & Z & \longleftarrow & Y \\
 \downarrow & & \downarrow & & \downarrow \\
 GX & \longrightarrow & GZ & \longleftarrow & GY
 \end{array}$$



Process Semantics

Linear Time-Branching Time Spectrum

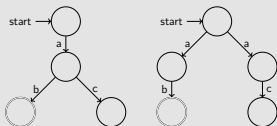
van Glabbeek '90

Graded Semantics

Coalgebraic framework for logics

Milius et al. '15

$$X \xrightarrow{\gamma} GX \xrightarrow{\alpha} M_1 X$$



Automaton Theory

Behavioural equivalence

\neq

Language semantics

Generalized Powerset Construction

EM-Law

Eilenberg-Moore distributive law $\zeta: TF \Rightarrow FT$, natural transformation compatible with the structure of T

EM-Semantics

Determinization of coalgebra $\gamma: X \rightarrow FTX$:

$$\gamma^\# : TX \xrightarrow{T\gamma} TFTX \xrightarrow{\zeta^{TX}} FTTX \xrightarrow{F\mu_X} FTX$$

Behavioural equivalence in $\gamma^\# \approx$ Language semantics
(nondeterministic, weighted, probabilistic automata)

Graded Monads

Graded Monads \mathbb{M}

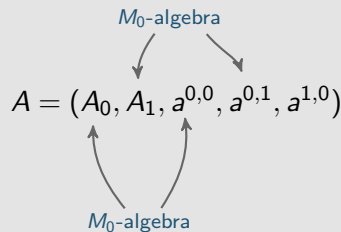
- Functors $M_n: \mathbf{C} \rightarrow \mathbf{C}$ for $n \in \mathbb{N}$
- Multiplications $\mu^{ij}: M_i M_j \Rightarrow M_{i+j}$
- Unit $\eta: Id \Rightarrow M_0$

+ monad laws (with indices)

M_n -Algebras

- Carriers A_k for $k \leq n$
- Structures $a^{ij}: M_i A_j \Rightarrow A_{i+j}$

+ algebra laws (with indices)



EM-Laws Define Graded Monads

Graded Monad \mathbb{M}_G

- $M_n = G^n$
- All natural transformations identity

Captures (finite-depth) Behavioural equivalence.

Graded Monad \mathbb{M}_ζ

- $M_n = F^n T$
- Multiplications defined by ζ
- η inherited from T

Captures *EM*-Semantics.

Graded Semantics

Definition

Graded semantics $(\mathbb{M}, \alpha: G \Rightarrow M_1)$

For $\gamma: X \rightarrow GX$ define inductively $\gamma^{(k)}: X \rightarrow M_k \mathbf{1}$:

$$\gamma^{(0)}: X \xrightarrow{\eta} M_0 X \xrightarrow{M_0!} M_0 \mathbf{1}$$

$$\gamma^{(k+1)}: X \xrightarrow{\alpha \cdot \gamma} M_1 X \xrightarrow{M_1 \gamma^{(k)}} M_1 M_k \mathbf{1} \xrightarrow{\mu^{1k}} M_{k+1} \mathbf{1}$$

Depth-1 Graded Monads

Definition

\mathbb{M} is *depth-1* if the following diagram is a coequalizer:

$$M_1 M_0 M_0 \begin{array}{c} \xrightarrow{M_1 \mu^{00}} \\ \xrightarrow{\mu^{10} M_0} \end{array} M_1 M_0 \xrightarrow{\mu^{10}} M_1$$

Equivalent: Depth-1 algebraic theory

Lemma

If \mathbb{M} depth-1

$\Rightarrow (M_n \mathbf{1}, M_{n+1} \mathbf{1}, \mu^{0,n}, \mu^{0,n+1}, \mu^{1,n})$ is free over its 0-part.

Coalgebraic Logic

Constants
Modalities

Syntactic Components given by $\mathcal{L} = (\Theta, \mathcal{O}, \Lambda)$

Propositional Operators

- $\theta \in \Theta$ $\hat{\theta}: 1 \rightarrow \Omega$
- $p \in \mathcal{O}$ $\llbracket p \rrbracket: \Omega^n \rightarrow \Omega$
- $\lambda \in \Lambda$ $\llbracket \lambda \rrbracket: G\Omega \rightarrow \Omega$

Semantics for a coalgebra $\gamma: X \rightarrow GX$

$$\llbracket \lambda \phi \rrbracket_{\gamma}: X \xrightarrow{\gamma} GX \xrightarrow{G\llbracket \phi \rrbracket_{\gamma}} G\Omega \xrightarrow{\llbracket \lambda \rrbracket} \Omega$$

Graded Logics

Definition

\mathcal{L} is a *graded logic* if

- Ω carries an M_0 -algebra (Ω, o)
- $\llbracket p \rrbracket : \Omega^n \rightarrow \Omega$ algebra homomorphism
- $\llbracket \lambda \rrbracket = f \circ \alpha$ $(\Omega, \Omega, o, o, f)$ is an M_1 -algebra

Theorem

\mathbb{M} depth-1 graded monad, \mathcal{L} graded logic

$\Rightarrow \mathcal{L}$ invariant for (α, \mathbb{M})

Characterizing Modalities for \mathbb{M}_ζ

$$\begin{array}{ccccc}
 FTT\Omega & \xrightarrow[\text{FTo}]{F\mu_\Omega} & FT\Omega & \xrightarrow{[\lambda]} & \Omega \\
 & & \searrow \text{Fo} & & \nearrow h \\
 & & & & F\Omega
 \end{array}$$

Theorem

Graded modal operators $[\lambda]: FT\Omega \rightarrow \Omega$ correspond to algebra-homomorphisms $h: \tilde{F}(\Omega, o) \rightarrow (\Omega, o)$.

Expressivity

Depth-0 separation

$$\hat{\theta}^*: M_0 1 \rightarrow \Omega$$

jointly injective

Depth-1 separation

$\forall A$ canonical (free over $(-)_0$)

$\forall \Phi \subseteq A_0 \rightarrow \Omega$ jointly injective, closed under \mathcal{O}

$\Rightarrow \llbracket \lambda \rrbracket(\llbracket \phi \rrbracket)$ jointly injective

Theorem

If \mathcal{L} is depth-0 and depth-1 separating, then \mathcal{L} is expressive.

$$\begin{array}{ccc}
 & & \llbracket \phi \rrbracket \in \Phi \\
 & & \swarrow \\
 M_1 A_0 & \xrightarrow{M_1 \llbracket \phi \rrbracket} & M_1 \Omega \\
 \downarrow a^{10} & & \downarrow f \\
 A_1 & \xrightarrow{\llbracket \lambda \rrbracket(\llbracket \phi \rrbracket)} & \Omega
 \end{array}
 \quad \leftarrow f \circ \alpha \in \Lambda$$

Reducing from Linear Time to Branching Time

$\mathcal{L} = (\Theta, \mathcal{O}, \Lambda)$ graded logic for $(\text{id}, \mathbb{M}_\zeta)$ on FT -coalgebras

\Downarrow

$$\Lambda' = \{h: F\Omega \rightarrow \Omega \mid h \circ F\omega \in \Lambda\}$$

$\mathcal{L}' = (\Theta, \mathcal{O}, \Lambda')$ graded logic for $(\text{id}, \mathbb{M}_F)$ on F -coalgebras

Proposition

If \mathcal{L}' is depth-1 separating, then \mathcal{L} is depth-1 separating,

Example: Machine Functor

T arbitrary, $F = 2 \times (-)^\Sigma$, $\Omega = 2$

Λ' : $\langle \top \rangle$ first projection, $\langle \sigma \rangle$ second projection (for $\sigma \in \Sigma$)

- Morphisms in **EM**(T) ✓
 \Rightarrow Modal operators in Λ are invariant
- Depth-1 separating for \mathbb{M}_F ✓
 $\Rightarrow (\emptyset, \emptyset, \Lambda)$ is depth-1 separating

\mathcal{V} -enriched Categories

Quantale

$$\mathcal{V} = (V, \leq, \otimes, k)$$

(V, \leq) complete lattice

(V, \otimes, k) monoid

$$(\bigvee_I u_i) \otimes v = \bigvee_I (u_i \otimes v)$$

Example: **Met**

$M_n 1$ metric spaces,
 $\llbracket \phi \rrbracket : X \rightarrow [0, 1]$

Instances

- (hemi-/pseudo-)metric spaces
- preorders (posets/setoids/sets)

Conclusion

What We Showed

- Graded semantics can capture EM-Semantics
- Invariance and expressivity follow from general results
- Quantalic generality

Future Work

- Fixpoints \Rightarrow relation to regular expressions
- Quantitative Kleisli semantics